

Note: Unless otherwise indicated, you need not show any work.

1. Both of the following are *group* codes. For each one: **(i)** determine the minimum nonzero weight of the code words then, using this value, determine **(ii)** the maximum number of errors that can reliably be detected and **(iii)** the maximum number of errors that can reliably be corrected.

(a) $\mathcal{C}_2 = \{(000\ 000\ 0000), (100\ 100\ 1011), (010\ 010\ 1101), (001\ 001\ 1110),$
 $(110\ 110\ 0110), (101\ 101\ 0101), (011\ 011\ 0011), (111\ 111\ 1000)\}$ in \mathbb{Z}_2^{10}

Ans: (i) The nonzero weights are 5, 6 and 7. The **minimum is 5**. This is the minimum distance so: (ii) **4 errors** can always be detected. (iii) **2 errors** can always be corrected.

(b) $\mathcal{C}_1 = \{(000\ 0000), (100\ 1011), (010\ 0111), (001\ 1101),$
 $(110\ 1100), (101\ 0110), (011\ 1010), (111\ 0001)\}$ in \mathbb{Z}_2^7

Ans: (i) The **minimum nonzero weight is 4**. This is the minimum distance so: (ii) **3 errors** can always be detected. (iii) **1 error** can always be corrected.

2. The following is a generator matrix for a group code: $G = \left(\begin{array}{cccc|ccc} 1 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 \end{array} \right)$

Do the following.

- (a) Find the code words corresponding to the messages $v = (1100)$ and $w = (1001)$.

Ans: $vG = (1100011)$ and $wG = (1001101)$

- (b) Write out the parity check matrix H corresponding to G .

Ans: $H = \left(\begin{array}{cccc|ccc} 1 & 1 & 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 \end{array} \right)$

- (c) For the received words (i) $r = (1000\ 001)$ and (ii) $s = (1000\ 100)$ find Hr^{tr} and Hs^{tr} , then use those results to correct both if there is an error. If you decide this can't be done, give a reason.

Ans: $Hr^{\text{tr}} = (111)^{\text{tr}}$, the 3rd column of H , so change the 3rd bit of r to get **(1010 001)**
 $Hs^{\text{tr}} = (010)^{\text{tr}}$, the 6th column of H , so change the 6th bit of s to get **(1000 110)**.

- (d) If possible, decode the corrected words in part (c) to obtain their original messages.

Ans: For r : **(1010)**, and for s : **(1000)**.