

Note: Unless otherwise indicated, you need not show any work.

1. Both of the following are *group* codes. For each one: **(i)** determine the minimum nonzero weight of the code words then, using this value, determine **(ii)** the maximum number of errors that can reliably be detected and **(iii)** the maximum number of errors that can reliably be corrected.

(a)  $\mathcal{C}_1 = \{(000\ 0000), (100\ 1011), (010\ 0111), (001\ 1101), (110\ 1100), (101\ 0110), (011\ 1010), (111\ 0001)\}$  in  $\mathbb{Z}_2^7$

**Ans:** (i) The **minimum nonzero weight is 4**. This is the minimum distance so: (ii) **3 errors** can always be detected. (iii) **1 error** can always be corrected.

(b)  $\mathcal{C}_2 = \{(000\ 000\ 0000), (100\ 101\ 1011), (010\ 111\ 1101), (001\ 110\ 1110), (110\ 010\ 0110), (101\ 011\ 0101), (011\ 001\ 0011), (111\ 100\ 1000)\}$  in  $\mathbb{Z}_2^{10}$

**Ans:** (i) The nonzero weights are 5, 6 and 7. The **minimum is 5**. This is the minimum distance so: (ii) **4 errors** can always be detected. (iii) **2 errors** can always be corrected.

2. The following is a generator matrix for a group code:  $G = \left( \begin{array}{cccc|cccc} 1 & 0 & 0 & 0 & 1 & 0 & 1 & \\ 0 & 1 & 0 & 0 & 1 & 1 & 1 & \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 & \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 & \end{array} \right)$

Do the following.

- (a) Find the code words corresponding to the messages  $v = (1100)$  and  $w = (0110)$ . Don't show your work. **Ans:**  $vG = (1100010)$  and  $wG = (0110100)$

- (b) Write out the parity check matrix  $H$  corresponding to  $G$ .

**Ans:**  $H = \left( \begin{array}{cccc|cccc} 1 & 1 & 0 & 1 & 1 & 0 & 0 & \\ 0 & 1 & 1 & 1 & 0 & 1 & 0 & \\ 1 & 1 & 1 & 0 & 0 & 0 & 1 & \end{array} \right)$

- (c) For the received words (i)  $r = (1000\ 001)$  and (ii)  $s = (1111\ 000)$  find  $Hr^{\text{tr}}$  and  $Hs^{\text{tr}}$ , then use those results to correct both if there is an error. If you decide this can't be done, give a reason.

**Ans:**  $Hr^{\text{tr}} = (100)^{\text{tr}}$ , the 5th column of  $H$ , so change the 5th bit to get **(1000101)**  
 $Hs^{\text{tr}} = (111)^{\text{tr}}$ , the 2nd column of  $H$ , so change the 2nd bit to get **(1011000)**.

- (d) If possible, decode the corrected words in part (c) to obtain their original messages.

**Ans:** For  $r$ : **(1000)**, and for  $s$ : **(1011)**.